

# Number Theory & Cryptography - Weekly meetings

Magma Tutorial [Part II]



Luis J. Dominguez Perez  
CIMAT-Zac, Enero 28 de 2015

# Surprise hands-on

Code the extended Euclidean algorithm

# Contenido, sección I

Polynomials

Finite fields

Primitive Elements

Plotting

# Rings

## Supported types of rings

- IntegerRing(.)
- Integers()
- Rationals()
- RealField()

```
Z11:=IntegerRing(11);  
w:=PrimitiveElement(Z11); w;  
[w*i:i in [1..10]];  
[w^i:i in [1..10]];  
Order(w);  
Category(w);  
Parent(w);
```

# Polynomial Rings

```
Z:=Integers();  
S:=PolynomialRing(Z);  
AssignNames(~S, ["x"]);
```

```
S<x>:=PolynomialRing(Z);
```

```
S<x,y>:=PolynomialRing(Integers(),2);
```

```
Z<x>:=PolynomialRing(Rationals());
```

# Hands-on

Define the following polynomials:

```
Z<x>:=PolynomialRing(Integers());  
px:=36*x^4+36*x^3+24*x^2+6*x+1;  
z:=2^64;  
p:=Evaluate(px,z);  
rx:=x^8-x^4+1;  
r:=Evaluate(rx,z);
```

The exercise:

- Write a short program to find a  $z$  such that  $p$  is a prime of 256 bits.
- Modify your program to find  $r$  *almost prime* of 256 bits.

# Contenido, sección 2

Polynomials

Finite fields

Primitive Elements

Plotting

# Finite fields

- `FiniteField()`
- `GaloisField()`
- `GF()`

If the field size is a prime power, then it is preferred to use `FiniteField`.

Extension field (not yet covered in the Arithmetic Module)

```
ext<F | n>
```

```
ExtensionField< F, x | polynomial>
```

# Elements in $\mathbb{F}_p$

## Example usage of elements in a field

```
p:=101;  
Fp:=FiniteField(p);  
a:=Random(p);  
b:=Fp!103;  
c:=Fp!-1000;  
(a+b) in Fp;  
(b*c) in Fp;  
a^(1999);  
Fp!a^(1999);  
a^(1999) eq Fp!a^(1999);
```

# Other functions for finite fields

Try these functions over  $F_p$

- `One(Fp);`
- `Identity(Fp);`
- `Zero(Fp);`
- `Fp.I;`

Now, create a quadratic extension of  $F_p$  and try the same functions. Then:

```
a2:=Random(Fp2); a2;  
myseq:=Eltseq(a2); myseq;  
myseq[1]*1+myseq[2]*Fp2.1;
```

# Hands-on

## Finite fields

- Generate a Finite field of 256 bits
- Generate a random element
- Find its inverse using Fermat's little theorem

# Bonus

```
> p:=RandomPrime(30);
> p;
471164429
> Factorization(p-1);
[ <2, 2>, <7, 1>, <16827301, 1> ]
> K:=GF(p);
> m:=Random(K)^7;
> m;
264249637
> P<x>:=PolynomialRing(K);
> f:=x^7-m;
> Factorization(f);
[
  <x + 168185695, 1>,
  <x + 266248969, 1>,
  <x + 301801881, 1>,
  <x + 331419936, 1>,
  <x + 426819080, 1>,
  <x + 427671341, 1>,
  <x + 433675243, 1>
]
> fact:=$1;
```

```
> for i:=1 to 7 do
for> z := -Evaluate (fact[i][1], 0);
for> print z, "^7 -", m, "=", z^7 - m;
for> end for;
302978734 ^7 - 264249637 = 0
204915460 ^7 - 264249637 = 0
169362548 ^7 - 264249637 = 0
139744493 ^7 - 264249637 = 0
44345349 ^7 - 264249637 = 0
43493088 ^7 - 264249637 = 0
37489186 ^7 - 264249637 = 0
>
```

# Contenido, sección 3

Polynomials

Finite fields

Primitive Elements

Plotting

# Primitive element

Let  $g$  be a multiplicative group of order  $n$ , and  $g \in G$ . The order of  $g$  divides  $n$ .

If  $b \in \mathbb{Z}_p^*$ , then  $b^{\Phi(n)} \equiv 1 \pmod{n}$ ,  $\#\mathbb{Z}_n^* = \Phi(n)$ .

Let  $p$ -prime,  $b \in \mathbb{Z}_p$ , then  $b^p \equiv b \pmod{p}$ , and  $\Phi(p) = p - 1$ . Also,  $b \not\equiv 0 \pmod{p}$  and  $\#\mathbb{Z}_p^* = p - 1$

Then,  $\exists \alpha \in \mathbb{Z}_p^*$  s.t.  $|\alpha| = p - 1$ , such an  $\alpha$  is called a *primitive element*.

# Hands on

## Exercise

- Make a program to find the primitive elements of a very small finite field, and verify the order of the generated group with your Euler totient function.

# Finding primitive elements in $\mathbb{Z}_p^*$

An easy way to determine if a random element  $\alpha \in \mathbb{Z}_p^*$  is primitive arises when the factorization (in prime powers) of  $(p - 1)$  is known.

Let  $p = p_1^{e_1} \cdots p_k^{e_k}$

An element  $\alpha \in \mathbb{Z}_p^*$  is primitive iff

$$\alpha^{(p-1)/p_j} \not\equiv 1 \pmod{p}.$$

# Finding primitive elements in $\mathbb{Z}_p^*$ II

## Proof

Let  $d = |\alpha|$ . We know  $d \mid (p-1)$ , and  $\alpha$  is primitive iff  $d = p-1$ .

Suppose  $\alpha^{(p-1)/p_j} \equiv 1 \pmod{p}$  for some  $j$ , then  $d \leq (p-1)/p_j$ , so  $d \neq (p-1)$ .

Now, suppose  $\alpha^{(p-1)/p_j} \not\equiv 1 \pmod{p}$  for  $1 \leq j \leq k$ . Suppose  $d \neq p-1$ , since  $d$  is a divisor of  $p-1$ , and  $d < p-1$ ,  $\exists p_j (1 \leq j \leq k)$  s.t.  $p_j$  is a divisor of  $(p-1)/d$ , but this implies  $d$  is a divisor of  $(p-1)/p_j$ .

Hence,  $\alpha^{(p-1)/p_j} \equiv \alpha^d \equiv 1 \pmod{p}$ , which is a contradiction  $\square$

# Finding primitive elements in $\mathbb{Z}_p^*$ III

One way to find a primitive element is using *Las Vegas* algorithm: choosing random  $\alpha$  and testing them, until a primitive element is found.

There are exactly  $\Phi(p-1)$  primitive elements in  $\mathbb{Z}_p^*$ , the probability to find a random primitive element is  $\Phi(p-1)/(p-1)$ .

Suppose  $p$  and  $p_1$  are prime, and  $p = 2p_1 + 1$ . Suppose  $\alpha \in \mathbb{Z}_p^*$  and  $\alpha \not\equiv 1 \pmod{p}$ , then  $\alpha$  is a primitive element iff  $\alpha^{(p-1)/2} \not\equiv 1 \pmod{p}$ .

# Finding primitive elements in $\mathbb{Z}_p^*$ IV

## *Proof*

Observe that  $\alpha^{(p-1)/p_j} \equiv \alpha^2 \pmod{p}$ , and  $\alpha^2 \equiv 1 \pmod{p}$  iff  $\alpha \equiv \pm 1 \pmod{p}$ . Follow the previous slides.  $\square$

## *Following*

If  $\alpha \not\equiv 1 \pmod{p}$ , and  $\alpha$  is not primitive, then  $\alpha^{(p-1)/2} \equiv 1 \pmod{p}$ , but then:

$$\begin{aligned} (-\alpha)^{(p-1)/2} &\equiv (-1)^{(p-1)/2} \alpha^{(p-1)/2} \pmod{p} \\ &\equiv (-1)^{(p-1)/2} \pmod{p} \\ &\equiv -1 \pmod{p} \end{aligned}$$

Then,  $-\alpha$  must be primitive.

# Contenido, sección 4

Polynomials

Finite fields

Primitive Elements

Plotting

# Gnuplot

In your research, you may end up needing to plot your results, either from Magma, Maple, C/C++ or Java.

Some languages have support for graphs, it is not the case of Magma, but we can use Gnuplot.

Gnuplot is a portable command-line driven graphing utility for Linux, Windows, and some other platforms. Other options are available.

# Hands on

To plot, first we need data.

- Generate 3 finite fields of 80, 112, and 128 bits
- Generate 3 random elements in the field
- time the exponentiation for several exponents
- Generate a table with the results and output to a file

# Setting up

```
set border 3
```

```
set terminal png picsize 512 512
```

```
set output "myfile.png"
```

```
set term postscript eps enhanced color
```

```
set output "myfile.eps"
```

At the end, one can convert the `epstopdf myfile.eps` to get an embeddable pdf for your article.

# Setting up

```
set style line 1 linetype 1 linewidth 1 \  
    pointtype 9 linecolor rgb "red"  
set style dots...  
set style point...
```

```
set key right bottom box
```

```
set title "My graph"  
set xlabel "Equivalent AES security level"  
set xrange [80:128]  
set xtics 88,96,104,112,120,128  
set mxtics 4
```

# Plotting

```
plot "Example" using 1:3 title 'Data1' with \
linespoints linestyle 1,\
    "Example" using 1:4 title 'Data2' with \
    linespoints linestyle 2,\
    "Example" using 1:5 title 'Data3' with \
    linespoints linestyle 3
```

Don't forget to reset your settings at the end of your plot!

# Multiplot

```
set size 1,1  
set origin 0,0  
set multiplot  
set xlabel "Security level"
```

```
set size 0.5,0.5
```

```
set origin 0,0.5  
set ylabel "CPU cycles"  
plot ...
```

```
set origin 0.5,0.5  
set ylabel "Milli seconds"  
plot ...
```

# More plot types

If bored, have a look at...

http:

[`//www.phyast.pitt.edu/~zov1/gnuplot/html/intro.html`](http://www.phyast.pitt.edu/~zov1/gnuplot/html/intro.html)<sup>1</sup>

---

<sup>1</sup> Obviously, you need some data to plot...

# End

## End of Part II

There's no part III